Safe Automated Refactoring for Intelligent Parallelization of Java 8 Streams

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***Abstract*—Streaming APIs are becoming more pervasive in mainstream Object-Oriented programming languages. For example, the Stream API introduced in Java 8 allows for functional-like, MapReduce-style operations in processing both finite and infinite data structures. However, using this API efficiently involves subtle considerations like determining when it is best for stream operations to run in parallel, when running operations in parallel can be less efficient, and when it is safe to run in parallel due to possible lambda expression side-effects. In this paper, we present an automated refactoring approach that assists developers in writing efficient stream code in a semantics-preserving fashion. The approach, based on a novel data ordering and typestate analysis, consists of preconditions for automatically determining when it is safe and possibly advantageous to convert sequential streams to parallel and unorder or de-parallelize already parallel streams. The approach was implemented as a plug-in to the Eclipse IDE, uses the WALA and SAFE analysis frameworks, and was evaluated on 11 Java projects consisting of ∼642K lines of code. We found that 57 of 157 candidate streams (36.31%) were refactorable, and an average speedup of 3.49 on performance tests was observed. The results indicate that the approach is useful in optimizing stream code to their full potential.**

***Index Terms*—refactoring, static analysis, automatic parallelization, typestate analysis, Java 8, streams**

I. Introduction

Streaming APIs are widely-available in today's mainstream, Object-Oriented programming languages and platforms [1], including Scala [2], JavaScript [3], C# [4], Java [5], and Android [6]. These APIs incorporate MapReduce-like [7] operations on native data structures such as collections. Below is a "sum of even squares" example in Java [1], where `map()` accepts a λ-expression (unit of computation) and results in the list element's square. The λ-expression argument to `filter()` evaluates to true iff the element is even:

```
list.stream().filter(x->x%2==0).map(x->x*x).sum();
```

MapReduce, which helps reduce the complexity of writing parallel programs by facilitating big data processing on multiple nodes using succinct functional-like programming constructs, is a popular programming paradigm for writing a specific class of parallel programs. It makes writing parallel code easier, as writing such code can be difficult due to possible data races, thread interference, and contention [8]–[10]. For instance, the code above can execute in parallel simply by replacing `stream()` with `parallelStream()`.

MapReduce, though, traditionally operates in a highly-distributed environment with no concept of shared memory, while Java 8 Stream processing operates in a single node under multiple threads or cores in a shared memory space. In the latter case, because the data structures for which the MapReduce-like operations execute are on the local machine, problems may arise from the close intimacy between shared memory and the operations being performed. Developers, thus, must *manually* determine whether running stream code in parallel results in an efficient yet interference-free program [11] and ensure that no operations on different threads interleave [12].

Despite the benefits [13, Ch. 1], using streams efficiently requires many subtle considerations. For example, it is often not straight-forward if running a particular operation in parallel is more optimal than running it sequentially due to potential side-effects of λ-expressions, buffering, etc. Other times, using *stateful* λ-expressions, i.e., those whose results depend on any state that may change during execution, can undermine performance due to possible thread contention. In fact, ∼4K stream questions have been posted on Stack Overflow [14], of which ∼5% remain unanswered, suggesting that there is developer confusion surrounding this topic.

In general, these kinds of errors can lead to programs that undermine concurrency, underperform, and are inefficient. Moreover, these problems may not be immediately evident to developers and may require complex interprocedural analysis, a thorough understanding of the intricacies of a particular stream implementation, and knowledge of situational API replacements. Manual analysis and/or refactoring (semantics-preserving, source-to-source transformation) to achieve optimal results can be overwhelming and error- and omission-prone. This problem is exacerbated by the fact that 157 total candidate streams[1] across 11 projects with a 34 project maximum[2] were found during our experiments (§ IV), a number that can increase over time as streams rise in popularity. In fact, Mazinanian *et al.* [15] found an increasing trend in the adoption of λ-expressions, an essential part of using the Java 8 stream API, with the number of λ-expressions being introduced increasing by two-fold between 2015 and 2016. And, a recent GitHub search by the authors yielded 350K classes importing the `java.util.stream` package.

The operations issued per stream may be many; we found an average of 4.14 operations per stream. Permutating through operation combinations and subsequently assessing performance, if such dedicated tests even exist, can be burdensome. (Manual)

[1]The number of candidate streams is affected by several analysis parameters, which involve performance trade-offs, as described in § IV-B and IV-C.

[2]A stream instance approximation is defined as an invocation to a stream API returning a stream object, e.g., `stream()`, `parallelStream()`.

interprocedural and type hierarchy analysis may be needed to discover ways to use streams in a particular context optimally.

Recently, attention has been given to retrofitting concurrency on to existing sequential (imperative) programs [16]–[18], translating imperative code to MapReduce [19], verifying and validating correctness of MapReduce-style programs [20]–[23], and improving performance of the underlying MapReduce framework implementation [24]–[27]. Little attention, though, has been paid to mainstream languages utilizing functional-style APIs that facilitate MapReduce-style operations over native data structures like collections. Furthermore, improving imperative-style MapReduce code that has either been handwritten or produced by one the approaches above has, to the best of our knowledge, not been thoroughly considered. Tang *et al.* [11] only briefly present preliminary progress towards this end, while Khatchadourian *et al.* [28] discuss engineering aspects.

The problem may also be handled by compilers or run times, however, refactoring has several benefits, including giving developers more control over where the optimizations take place and making parallel processing explicit. Refactorings can also be issued multiple times, e.g., prior to major releases, and, unlike static checkers, refactorings transform source code, a task that can be otherwise error-prone and involve nuances.

We propose a fully-automated, semantics-preserving refactoring approach that transforms Java 8 stream code for improved performance. The approach is based on a novel data ordering and typestate analysis. The ordering analysis involves inferring when maintaining the order of a data sequence in a particular expression is necessary for semantics preservation. Typestate analysis is a program analysis that augments the type system with "state" and has been traditionally used for preventing resource errors [29], [30]. Here, it is used to identify stream usages that can benefit from "intelligent" parallelization, resulting in more efficient, semantically-equivalent code.

Typestate was chosen to track state changes of streams that may be aliased and to determine the final state following a terminal (reduction) operation. Non-terminal (intermediate) operations may return the receiver stream, in which case traditional typestate applies. However, we augmented typestate to apply when a *new* stream is returned in such situations (cf. § III-B and III-D). Our approach interprocedurally analyzes relationships between types. It also discovers possible side-effects in arbitrarily complex λ-expressions to safely transform streams to either execute sequentially or in parallel, depending on which refactoring preconditions, which we define, pass. Furthermore, to the best of our knowledge, it is the first automated refactoring technique to integrate typestate.

The refactoring approach was implemented as an open-source Eclipse [31] plug-in that integrates analyses from WALA [32] and SAFE [33]. The evaluation involved studying the effects of our plug-in on 11 Java projects of varying size and domain with a total of $\sim$642K lines of code. Our study indicates that (i) given its interprocedural nature, the (fully automated) analysis cost is reasonable, with an average running time of 0.45 minutes per candidate stream and 6.602 seconds per thousand lines of code, (ii) despite their ease-of-use, parallel streams are not commonly (manually) used in modern Java software, motivating an automated approach, and (iii) the proposed approach is useful in refactoring stream code for greater efficiency despite its conservative nature. This work makes the following contributions:

Precondition formulation. We present a novel refactoring approach for maximizing the efficiency of their Java 8 stream code by automatically determining when it is safe and possibly advantageous to execute streams in parallel, when running streams in parallel can be counterproductive, and when ordering is unnecessarily depriving streams of optimal performance. Our approach refactors streams for greater parallelism while maintaining original semantics.

Generalized typestate analysis. Streams necessitate several generalizations of typestate analysis, including determining object state at arbitrary points and support for immutable object call chains. Reflection is also combined with (hybrid) typestate analysis to identify initial states.

Implementation and experimental evaluation. To ensure real-world applicability, the approach was implemented as an Eclipse plug-in built on WALA and SAFE and was used to study 11 Java programs that use streams. Our technique successfully refactored 36.31% of candidate streams, and we observed an average speedup of 3.49 during performance testing. The experimentation also gives insights into how streams are used in real-world applications, which can motivate future language and/or API design. These results advance the state of the art in automated tool support for stream code to perform to their full potential.

II. Motivation, Background, and Problem Insight

We present a running example that highlights some of the challenges associated with analyzing and refactoring streams for greater parallelism and increased efficiency. Lst. 1 portrays code that uses the Java 8 Stream API to process collections of `Widget`s with `weight`s (class not shown). Lst. 1a is the original version, while Lst. 1b is the improved (but semantically-equivalent) version after our refactoring. In lst. 1a, a `Collection` of `Widget`s is declared (line 1) that does not maintain element ordering as `HashSet` does not support it [34]. Note that ordering is dependent on the run time type.

A `stream` (a view representing element sequences supporting MapReduce-style operations) of `unorderedWidgets` is created on line 5. It is sequential, meaning its operations will execute serially. Streams may also have an *encounter order*, which can be dependent on the stream's source. In this case, it will be unordered since `HashSet`s are unordered.

On line 6, the stream is `sorted` by the corresponding *intermediate* operation, the result of which is a (possibly) new stream with the encounter order rearranged accordingly. `Widget::getWeight` is a method *reference* denoting the method that should be used for the comparison. Intermediate operations are deferred until a *terminal* operation is executed like `collect()` (line 7). `collect()` is a special kind of (mutable) reduction that aggregates results of prior intermediate operations into a given `Collector`. In this case, it is one that yields a `List`. The result is a `Widget List` sorted by `weight`.

Listing 1 Snippet of `Widget` collection processing using Java 8 streams based on *java.util.stream (Java SE 9 & JDK 9)* [5].

(a) Stream code snippet prior to refactoring.

```
1  Collection<Widget> unorderedWidgets = new HashSet<>();
2  Collection<Widget> orderedWidgets = new ArrayList<>();
3
4  List<Widget> sortedWidgets = unorderedWidgets
5    .stream()
6    .sorted(Comparator.comparing(Widget::getWeight))
7    .collect(Collectors.toList());
8
9  // collect weights over 43.2 into a set in parallel.
10 Set<Double> heavyWidgetWeightSet = orderedWidgets
11   .parallelStream().map(Widget::getWeight)
12   .filter(w -> w > 43.2).collect(Collectors.toSet());
13
14 // sequentially collect into a list, skipping first 1000.
15 List<Widget> skippedWidgetList = orderedWidgets
16   .stream().skip(1000).collect(Collectors.toList());
17
18 // collect the first green widgets into a list.
19 List<Widget> firstGreenList = orderedWidgets
20   .stream()
21   .filter(w -> w.getColor() == Color.GREEN)
22   .unordered().limit(5).collect(Collectors.toList());
23
24 // collect distinct widget weights into a TreeSet.
25 Set<Double> distinctWeightSet = orderedWidgets
26   .stream().parallel()
27   .map(Widget::getWeight).distinct()
28   .collect(Collectors.toCollection(TreeSet::new));
29
30 // collect distinct widget colors into a HashSet.
31 Set<Color> distinctColorSet = orderedWidgets
32   .parallelStream().map(Widget::getColor)
33   .distinct()
34   .collect(HashSet::new, Set::add, Set::addAll);
```

(b) Improved stream client code via refactoring.

```
1  Collection<Widget> unorderedWidgets = new HashSet<>();
2  Collection<Widget> orderedWidgets = new ArrayList<>();
3
4  List<Widget> sortedWidgets = unorderedWidgets
5    .~~stream()~~parallelStream()
6    .sorted(Comparator.comparing(Widget::getWeight))
7    .collect(Collectors.toList());
8
9  // collect weights over 43.2 into a set in parallel.
10 Set<Double> heavyWidgetWeightSet = orderedWidgets
11   .parallelStream().map(Widget::getWeight)
12   .filter(w -> w > 43.2).collect(Collectors.toSet());
13
14 // sequentially collect into a list, skipping first 1000.
15 List<Widget> skippedWidgetList = orderedWidgets
16   .stream().skip(1000).collect(Collectors.toList());
17
18 // collect the first green widgets into a list.
19 List<Widget> firstGreenList = orderedWidgets
20   .~~stream()~~parallelStream()
21   .filter(w -> w.getColor() == Color.GREEN)
22   .unordered().limit(5).collect(Collectors.toList());
23
24 // collect distinct widget weights into a TreeSet.
25 Set<Double> distinctWeightSet = orderedWidgets
26   .stream().~~parallel()~~
27   .map(Widget::getWeight).distinct()
28   .collect(Collectors.toCollection(TreeSet::new));
29
30 // collect distinct widget colors into a HashSet.
31 Set<Color> distinctColorSet = orderedWidgets
32   .parallelStream().map(Widget::getColor)
33   .unordered().distinct()
34   .collect(HashSet::new, Set::add, Set::addAll);
```

It may be possible to increase performance by running this stream's "pipeline" (i.e., its sequence of operations) in parallel. Lst. 1b, line 5 displays the corresponding refactoring with the stream pipeline execution in parallel (removed code is ~~struck through~~, while the added code is underlined). Note, however, that had the stream been *ordered*, running the pipeline in parallel may result in worse performance due to the multiple passes and/or data buffering required by *stateful* intermediate operations (SIOs) like `sorted()`. Because the stream is *unordered*, the reduction can be done more efficiently as the framework can employ a divide-and-conquer strategy [5].

In contrast, line 2 instantiates an `ArrayList`, which maintains element ordering. Furthermore, a parallel stream is derived from this collection (line 11), with each `Widget` mapped to its weight, each weighted filtered (line 12), and the results `collect`ed into a `Set`. Unlike the previous example, however, no optimizations are available here as an SIO is not included in the pipeline and, as such, the parallel computation does not incur the aforementioned possible performance degradation.

Lines 15–16 create a list of `Widgets` gathered by (sequentially) skipping the first thousand from `orderedWidgets`. Like `sorted()`, `skip()` is also an SIO. Unlike the previous example, though, executing this pipeline in parallel could be counterproductive because, as it is derived from an *ordered* collection, the stream is ordered. It may be possible to unorder the stream (via `unordered()`) so that its pipeline would be more amenable to parallelization. In this situation, however, unordering could alter semantics as the data is assembled into a structure maintaining ordering. As such, the stream remains sequential as element ordering must be preserved.

On lines 19–22, the first five green `Widgets` of `orderedWidgets` are sequentially `collect`ed into a `List`. As `limit()` is an SIO, performing this computation in parallel could have adverse effects as the stream is ordered (with the source being `orderedWidgets`). Yet, on line 22, the stream is `unordered`[3] *before* the `limit()` operation. Because the SIO is applied to an unordered stream, to improve performance, the pipeline is refactored to parallel on line 20 in lst. 1b. Although similar to the refactoring on line 5, it demonstrates that stream ordering does not solely depend on its source.

A distinct widget weight `Set` is created on lines 25–28. Unlike the previous example, this collection *already* takes place in `parallel`. Note though that there is a possible performance degradation here as the SIO `distinct()` may require multiple passes, the computation takes place in *parallel*, *and* the stream is *ordered*. Keeping the parallel computation but unordering the stream may improve performance but we would need to determine whether doing so is safe, which can be error-prone if done manually, especially on large and complex projects.

Our insight is that, by analyzing the type of the resulting reduction, we may be able to determine if unordering a stream is safe. In this case, it is a (mutable) reduction (i.e., `collect()` on line 28) to a `Set`, of which subclasses that do not preserve ordering exist. If we could determine that the resulting `Set` is unordered, unordering the stream would be safe since the collection operation would not preserve ordering. The type of the resulting `Set` returned here is determined

[3]The use of `unordered()` is deliberate despite nondeterminism.

TABLE I
CONVERT SEQUENTIAL STREAM TO PARALLEL PRECONDITIONS.

	exe	ord	se	SIO	ROM	transformation
P1	seq	unord	*F*	N/A	N/A	Convert to para.
P2	seq	ord	*F*	*F*	N/A	Convert to para.
P3	seq	ord	*F*	*T*	*F*	Unorder and convert to para.

by the passed `Collector`, in this case, `Collectors.toCollection(TreeSet::new)`, the argument to which is a reference to the default constructor. Unfortunately, since `TreeSet`s preserve ordering, we must keep the stream ordered. Here, to improve performance, it may be advantageous to run this pipeline, perhaps surprisingly, *sequentially* (line 26, lst. 1b).

Lines 31–34 map, in parallel, each `Widget` to its `Color`, filter those that are `distinct`, and `collect`ing them into a `Set`. To demonstrate the variety of ways mutable reductions can occur, a more direct form of `collect()` is used rather than a `Collector`, and the collection is to a `HashSet`, which does *not* maintain element ordering. As such, though the stream is originally ordered, since the (mutable) reduction is to an *unordered* destination, we can infer that the stream can be safely unordered to improve performance. Thus, line 33 in lst. 1b shows the inserted call to `unordered()` immediately prior to `distinct()`. This allows `distinct()` to work more efficiently under parallel computation [5].

Manual analysis of stream client code can be complicated, even as seen in this simplified example. It necessitates a thorough understanding of the intricacies of the underlying computational model, a problem which can be compounded in more extensive programs. As streaming APIs become more pervasive, it would be extremely valuable to developers, particularly those not previously familiar with functional programming, if automation can assist them in writing efficient stream code.

III. OPTIMIZATION APPROACH

A. Intelligent Parallelization Refactorings

We propose two new refactorings, i.e., CONVERT SEQUENTIAL STREAM TO PARALLEL and OPTIMIZE PARALLEL STREAM. The first deals with determining if it is possibly advantageous (performance-wise, based on type analysis) and safe (e.g., no race conditions, semantics alterations) to transform a sequential stream to parallel. The second deals with a stream that is *already* parallel and ascertains the steps (transformations) necessary to possibly improve its performance, including *unordering* and *converting* the stream to sequential.

1) Converting Sequential Streams to Parallel: Table I portrays the preconditions for our proposed CONVERT SEQUENTIAL STREAM TO PARALLEL refactoring. It lists the conditions that must hold for the transformation to be both semantics-preserving as well as possibly advantageous, i.e., resulting in a possible performance gain. Column **exe** denotes the stream's execution mode, i.e., whether, upon the execution of a terminal operation, its associated pipeline will execute sequentially or in parallel ("seq" is sequential and "para" is parallel). Column **ord** denotes whether the stream is associated with an encounter order, i.e., whether elements of the stream *must* be visited in a particular order ("ord" is ordered and "unord" is unordered). Column **se** represents whether any behavioral parameters (λ-expressions) that will execute during the stream's pipeline have possible side-effects. Column **SIO** constitutes whether the pipeline has any stateful intermediate operations. Column **ROM** represents whether the encounter order must be preserved by the result of the terminal reduction operation. A *T* denotes that the reduction result depends on the encounter order of a previous (intermediate) operation. Conversely, an *F* signifies that any ordering of the input operation to the reduction need not be preserved. Column **transformation** characterizes the transformation actions to take when the corresponding precondition passes (note the conditions are mutually exclusive). *N/A* is either *T* or *F*.

TABLE II
OPTIMIZE PARALLEL STREAM PRECONDITIONS.

	exe	ord	SIO	ROM	transformation
P4	para	ord	*T*	*F*	Unorder.
P5	para	ord	*T*	*T*	Convert to seq.

A stream passing P1 is one that is sequential, unordered, and has no side-effects. Because this stream is already unordered, whether or not its pipeline contains an SIO is inconsequential. Since the stream is unordered, any SIOs can run efficiently in parallel. Moreover, preserving the ordering of the reduction is also inconsequential as no original ordering exists. Here, it is both safe and possibly advantageous to run the stream pipeline in parallel. The stream derived from `unorderedWidgets` on line 5, lst. 1 is an example of a stream passing P1. A stream passing P2 is also sequential and free of λ-expressions containing side-effects. However, such streams are ordered, meaning that the refactoring only takes place if no SIOs exist. P3, on the other hand, will allow such a refactoring to occur, i.e., if an SIO exists, *only* if the ordering of the reduction's result is inconsequential, i.e., the reduction ordering need not be maintained. As such, the stream can be *unordered* immediately before the (first) SIO (as performed on line 33, lst. 1b). The stream created on line 16, lst. 1 is an example of a stream failing this precondition.

2) Optimizing Parallel Streams: Table II depicts the preconditions for the OPTIMIZE PARALLEL STREAM refactoring. Here, the stream in question is *already* parallel. A stream passing either precondition is one that is ordered and whose pipeline contains an SIO. Streams passing P4 are ones where the reduction does not need to preserve the stream's encounter order, i.e., **ROM** is *F*. An example is depicted on line 32, lst. 1. Under these circumstances, the stream can be explicitly unordered immediately before the (first) SIO, as done on line 33 of lst. 1b. Streams passing P5, on the other hand, are ones that the reduction ordering *does* matter, e.g., the stream created on line 26. To possibly improve performance, such streams are transformed to *sequential* (line 26, lst. 1b).

B. Identifying Stream Creation

Identifying where in the code streams are created is imperative for several reasons. First, streams are typically derived from a source (e.g., a collection) and take on its characteristics (e.g., ordering). This is used in tracking stream attributes across their pipeline (§ III-C). Second, for streams passing preconditions, the creation site serves a significant role in the transformation.

There are several ways to create streams, including being derived from `Collection`s, being created from arrays (e.g., `Arrays.stream()`), and via static factory methods (e.g., `IntStream.range()`). Streams may also be directly created via constructors, but it is not typical of clients, which are our focus. We consider stream creation point approximations as any expression evaluating to a type implementing the `java.util.stream.BaseStream` interface, which is the top-level stream interface. We exclude, however, streams emanating from intermediate operations, i.e., instance methods whose receiver *and* return types implement the stream interface, as such methods are not likely to produce new streams but rather ones derived from the receiver but with different attributes.

C. Tracking Streams and Their Attributes

We discuss our approach to tracking streams and their attributes (i.e., state) using a series of labeled transition systems (LTSs). The LTSs are used in in the typestate analysis (§ III-D).

1) Execution Mode:

Definition 1. The LTS E is a tuple $E = (E_S, E_\Lambda, E_\rightarrow)$ where $E_S = \{\bot, seq, para\}$ is the set of states, E_Λ is a set of labels, and $E_\rightarrow$ is a set of labeled transitions.

The labels E_Λ corresponds to method calls that either create or transform the execution mode of streams. We denote the initial stream ("phantom") state as $\bot$. Different stream creation methods may transition the newly created stream to one that is either sequential or parallel. Transitions stemming from the $\bot$ state represent stream creation methods (§ III-B). As an example, the stream on line 5, lst. 1a would transition from $\bot$ to *seq*, while the stream at line 26 would transition from *seq* to *para* as a result of the corresponding call.

2) Ordering: Whether a stream has an encounter order depends on the stream source (run time) type and the intermediate operations. Certain sources (e.g., `List`) are intrinsically ordered, whereas others (e.g., `HashSet`) are not. Some intermediate operations (e.g., `sorted()`) may impose an encounter order on an otherwise unordered stream, and others may render an ordered stream unordered (e.g., `unordered()`). Further, some terminal operations may ignore encounter order (e.g., `forEach()`) while others (e.g., `forEachOrderer()`) abide by it [5].

Definition 2. The LTS O for tracking stream ordering is the tuple $O = (O_S, O_\Lambda, O_\rightarrow)$ where $O_S = \{\bot, ord, unord\}$ and other components are in line with definition 1.

For instance, the stream on line 5, lst. 1a would transition from $\bot$ to *unord* due to `HashSet.stream()`. Although the compile-time type of `unorderedWidgets` is `Collection` (line 1), we use an interprocedural type inference algorithm (explained next) to approximate `HashSet`. The stream at line 26 would transition from $\bot$ to *ord* state as a result of `orderedWidgets` having the type `ArrayList` (line 2).

a) Approximating Stream Source Types and Characteristics: The fact that stream ordering can depend on the run time type of its source necessitates that its type be approximated. For this, we use an interprocedural type inference algorithm via points-to analysis [35] that computes the possible run time types of the receiver from which the stream is created. Once the type is obtained, whether source types produce ordered or unordered streams is determined via reflection. While details are in § IV-A, briefly, the type is reflectively instantiated and its `Spliterator` [36] extracted. Then, stream characteristics, e.g., ordering, are queried [36]. This is enabled by the fact that collections and other types supporting streams do not typically change their ordering characteristics dynamically.

Using reflection in this way amounts to a kind of *hybrid* typestate analysis where initial states are determined via dynamic analysis. If reflection fails, e.g., an abstract type is inferred, the default is to ordered and sequential. This choice is safe considering that there is no net effect caused by our proposed transformations, thus preserving semantics. Furthermore, to prevent ambiguity in state transitions, it is required that each inferred type have the same attributes.

D. Tracking Stream Pipelines

Tracking stream pipelines is essential in determining satisfied preconditions. Pipelines can arbitrarily involve multiple methods and classes as well as be data-dependent (i.e., spanning multiple branches). In fact, during our evaluation (§ IV), we found many real-world examples that use streams interprocedurally.

Our automated refactoring approach involves developing a variant of typestate analysis [29], [30] to track stream pipelines and determine stream state when a terminal operation is issued. Typestate analysis is a program analysis that augments the type system with "state" information and has been traditionally used for prevention of program errors such as those related to resource usage. It works by assigning each variable an initial ($\bot$) state. Then, (mutating) method calls transition the object's state. States are represented by a lattice and possible transitions are represented by LTSs. If each method call sequence on the receiver does not eventually transition the object back to the $\bot$ state, the object may be left in a nonsensical state, indicating the potential presence of a bug.

Our typestate analysis makes use of a call graph, which is created via a k-CFA call graph construction algorithm [37], making our analysis both object and context sensitive (the context being the k-length call string). In other words, it adds context so that method calls to an object creation site (**`new`** operator) can be distinguished from one another [38, Ch. 3.6]. It is used here to consider client-side invocations of API calls as object creations. Setting $k = 1$ would not suffice as the analysis would not consider the client contexts as stream creations. As such, at least for streams, k must be $>= 2$. Although k is flexible in our approach, we use $k = 2$ as the default for streams and $k = 1$ elsewhere. § IV-B1 discusses how k was set during our experiments, as well as a heuristic to help guide developers in choosing a sufficient k.

We formulate a variant of typestate since operations like `sorted()` return (possibly) new streams derived from the receiver stream with their attributes altered. Definition 3 portrays the formalism capturing the concept of typestate analysis used in the remainder of this section. Several generalizations are made to extract typestate at a particular program point.

Definition 3 (Typestate Analysis)**.** Define $\mathit{TState}_{\mathit{LTS}}(i_s, \mathit{exp}) = S$ where LTS is a labeled transition system, i_s a stream instance, exp an expression, and to be the possible states S of i_s at exp according to *LTS*.

In definition 3, exp, an expression in the Abstract Syntax Tree (AST), is used to expose the internal details of the analysis. Typically, typestate is used to validate complete statement sequences. Regarding definition 3, this would be analogous to exp corresponding to a node associated with the last statement of the program. In our case, we are interested in typestates at particular program points; otherwise, we may not be able to depict typestate at the execution of the terminal operation accurately.

As an example, let i_s be the stream on line 5, lst. 1a and exp the expression `collect()` at line 7. Then, $\mathit{TState}_O(i_s, \texttt{collect(..)}) = \{\mathit{ord}\}$.

Traditional typestate analysis is used with (mutating) methods that alter object state. The Stream API, though, is written in an immutable style where each operation returns a stream reference that may refer to a new object. A naïve approach may involve tracking the typestates of the returned references from intermediate operations. Doing so, however, would produce an undesirable result as each stream object would be at the starting state.

§ III-C treats intermediate operations as being (perhaps void returning) methods that mutate the state of the receiver. This makes the formalism concise. However, in actuality, intermediate operations are *value returning methods*, returning a reference to the same (general) type as the receiver. As such, the style of this API is that of immutability, i.e., "manipulating" a stream involves creating a new stream based on an existing one. In such cases, the receiver is then considered *consumed*, i.e., any additional operations on the receiver would result in a run time exception, similar to linear type systems [39].

Our generalized typestate analysis works by tracking the state of stream instances as follows. For a given expression, the analysis yields a set of possible states for a given instance following the evaluation of the expression. Due to the API style, a typestate analysis that has a notion of instances that are *based on* other instances is needed. As such, we compute the typestate of individual streams and proceed to *merge* the typestates to obtain the *final* typestate after the expression of where a terminal operation consumes the stream. The final typestate is derived at this point because that is when all of the (queued) intermediate operations will execute. Moreover, the final typestate is a *set* due to dataflow analysis of possible branching.

1) Intermediate Streams: A stream is created via APIs calls stemming from the $\bot$ state as discussed in § III-C. Recall that intermediate operations may or may not also create streams based on the receiver. We coin such streams as *intermediate* streams as they are used to progress the computation to a final result. Moreover, intermediate streams cannot be instantiated alone; they must be based on (or derived from) existing ones. If an intermediate stream is derived from another intermediate stream, then, there must exist a chain of intermediate stream creations that starts at a non-intermediate stream. Due to conditional branching and polymorphism, there may be multiple such (possible) chains. Intermediate streams must be appropriately arranged so that the correct final state may be computed.

To sequence stream instances, we require a "predecessor" function $\mathit{Pred}(i_s) = \{i_{s_1}, \ldots, i_{s_n}\}$ that maps a stream i_s to a set of streams that may have been used to create i_s. $\mathit{Pred}(i_s)$ is computed by using the points-to set of the reference used as the receiver when i_s was instantiated. Definition 4 describes this function more generally.

Definition 4 (Predecessor Objects)**.** Define $\mathit{Pred}(\texttt{o.m()}) = \{i_1, i_2, \ldots, i_n\}$ where `o` is an object reference, `m` a method, `o.m()` results in an object reference, and $i_k \in \{i_1, i_2, \ldots, i_n\}$ for $1 \leq k \leq n$ an abstract heap object identifier:

$$\mathit{Pred}(\texttt{o.m()}) = \begin{cases} \varnothing & \text{if } \texttt{m()} \text{ is not intermediate.} \\ \mathit{PointsTo}(\texttt{o}) & \text{o.w.} \end{cases}$$

2) Typestate Merging: Since intermediate operations possibly create new streams based on the receiver, the typestate analysis will generate different states for any stream produced by an intermediate operation. We are interested in, however, the final state just before the commencement of the terminal operation, which results in stream consumption. Recall from § III-C1 that $\bot$ models an initial state. As such, $\bot$ will symbolize the initial state of intermediate streams. In other words, although an intermediate stream may "inherit" state from the stream from which it is derived, in our formalism, we use $\bot$ as a placeholder until we can derive what exactly the state should be. To this end, we introduce the concept of *typestate merging*.

First, we define a state selection function that results in the first state if it is not $\bot$ and the second state otherwise:

Definition 5 (State Selection)**.** Define $\mathit{Select}\colon\ S \times S \to\ S$ to be the state selection function:

$$\mathit{Select}(s_i, s_j) = \begin{cases} s_j & \text{if } s_i = \bot \\ s_i & \text{o.w.} \end{cases}$$

Definition 5 "selects" the "most recent" state in the case that the typestate analysis determines it for the instance under question and a previous state otherwise. For example, let $s_i = \bot$ and $s_j = \mathit{para}$. Then, $\mathit{Select}(s_i, s_j) = \mathit{para}$. Likewise, let $s_i = \mathit{unord}$ and $s_j = \mathit{ord}$. Then, $\mathit{Select}(s_i, s_j) = \mathit{unord}$.

Next, we define the state merging function, which allows us to merge two sets of states, as follows:

Definition 6 (State Merging)**.** Define $\mathit{Merge}(S_i, S_j) = S$ to be the typestate merging function:

$$\mathit{Merge}(S_i, S_j) = \begin{cases} S_i & \text{if } S_j = \varnothing \\ S_j & \text{if } S_i = \varnothing \\ \{\mathit{Select}(s_i, s_j) \mid s_i \in S_i \land s_j \in S_j\} & \text{o.w.} \end{cases}$$

As an example, let $S_i = \{\bot\}$ and $S_j = \{\mathit{seq}, \mathit{para}\}$. Then, $\mathit{Merge}(S_i, S_j) = \{\mathit{seq}, \mathit{para}\}$. Likewise, let $S_i = \{\mathit{ord}, \mathit{unord}\}$ and $S_j = \{\mathit{ord}, \mathit{unord}\}$. Then, $\mathit{Merge}(S_i, S_j) = \{\mathit{unord}, \mathit{ord}\}$.

Finally, we define the notation of merged typestate analysis:

Definition 7 (Merged Typestate Analysis)**.** Define $\mathit{MTState}_{\mathit{LTS}}(i_s, \mathit{exp}) = S$ where LTS is a labeled transition system, i_s a stream, exp an expression, to be the typestate analysis merging function:

$$MTState_{LTS}(i_s, exp) = \begin{cases} TState_{LTS}(i_s, exp) & \text{if } Pred(\texttt{o.m()}) = \varnothing \\ \bigcup\limits_{i_{s_k} \in Pred(i_s)} Merge(TState_{LTS}(i_s, exp), \\ \qquad MTState_{LTS}(i_{s_k}, exp)) & \text{o.w.} \end{cases}$$

This final function aggregates typestate over the complete method call chain until the terminal operation after exp.

E. Determining Whether Reduction Ordering Matters

To obtain a result from stream computations, a terminal (reduction) operation must be issued. Determining whether the ordering of the stream immediately before the reduction matters (ROM) equates to discovering whether the reduction result is the same regardless of whether the stream is ordered or not. In other words, the result of the terminal operation does not depend on the ordering of the stream for which the operation is invoked, i.e., the value when the stream is ordered is equal to the value when the stream is unordered. Some reductions (terminal operations) do not return a value, i.e., they are void returning methods. In these cases, the *behavior* rather than the resulting value should be the same. Terminal operations fall into two categories, namely, those that produce a result, e.g., `count()`, and those that produce a side-effect, normally by accepting a λ-expression, e.g., `forEach()` [5].

1) Non-scalar Result Producing Terminal Operations: In the case of non-scalar return values, whether the return type maintains ordering is determined by reusing the reflection technique described in § III-C2a. Specifically, a stream is reflectively derived from an instance of the non-scalar return (run time) type approximations and its characteristics examined. And, from this, whether reduction order matters is determined as follows. If it is impossible for the returned non-scalar type to maintain an element ordering, e.g., it is a `HashSet`, then, the result ordering cannot make a difference in the program's behavior. If, on the other hand, the returned type *can* maintain an ordering, we conservatively determine that the reduction ordering *does* matter. As before, if there is any inconsistencies between the ordering characteristics of the approximated types, the default is ordered.

2) Side-effect Producing Terminal Operations: When there is a void return value, as is the case with side-effect producing terminal operations, then, we need to know the order in which the stream elements are "served" to the λ-expression argument producing the side-effect. Currently, the list of void terminal operations that maintain element ordering is also a parameter to our analysis. As with determining SIOs, a more sophisticated analysis would be needed to possibly approximate this characteristic. In the current Java 8 Stream API, there are only two such methods, namely, `forEach()` and `forEachOrdered()`.

3) Scalar Result Producing Terminal Operations: The last case is perhaps the most difficult. While discussing whether non-scalar types (e.g., containers) maintain element ordering seems natural, when the reduction is to a scalar type, it is challenging to determine whether or not the element ordering used to produce the resulting value had any influence over it. Another view of the problem involves determining whether or not the operation(s) "building" the result from the stream are associative. Examples of associative operations include numeric addition, minimum, and maximum, and string concatenation [5]. To address this, we divide the problem into determining the associativity of specialized and general reduction operations.

a) Specialized Reduction Operations: Luckily, the number and associativity property of specialized reduction operations are fixed. As such, the list of specialized operations along with their associativity property is input to the approach.

b) General Reduction Operations: The remaining general reduction operations are `reduce()` and `collect()`. We have already covered the cases where these operations return non-scalar types. What remains is the cases when these operations return scalar types. Due to the essence of `collect()`, in practice, the result type will most likely fall into the non-scalar category. In fact, `collect()` is a specialization of `reduce()` meant for mutable reductions. Recall from § II that such operations collect results in a container such as a collection [5].

The generality of these reduction operations make determining whether ordering matters difficult. For example, even a simple sum reduction can be difficult for an automated approach to analyze. Consider the following code [5] that adds `Widget` weights together using `reduce()`:

```
widgets.stream().reduce(0,
    (sum, b) -> sum + b.getWeight(), Integer::sum);
```

The first argument is the *identity* element; the second an *accumulator* function, adding a `Widget`'s weight into the accumulated sum. The last argument combines two integer sums by adding them. The question is how, in general, can we tell that this is performing an operation that is associative like summation? In other words, how can we determine that the reducer computation is independent of the order of its inputs? It turns out that this is precisely the *reducer commutativity* problem [20]. Unfortunately, this problem has been shown to be undecidable by Chen *et al.* [20]. While we will consider approximations and/or heuristics as future work, currently, our approach conservatively fails preconditions in this case. During our experiments detailed in § IV, these failures only accounted for 5%.

IV. Evaluation

A. Implementation

Our approach was implemented as a publicly available, open source Eclipse IDE [31] plug-in [28] and built upon WALA [32] and SAFE [33]. Eclipse is leveraged for its extensive refactoring support [40] and that it is completely open-source for all Java development. WALA is used for static analyses such as side-effect analysis (ModRef), and SAFE, which depends on WALA, for its typestate analysis. SAFE was altered for programmatic use and "intermediate" typestates (cf. § III-D2). For the refactoring portion, Eclipse ASTs with source symbol bindings are used as an intermediate representation (IR), while the static analysis consumes a Static Single Assignment (SSA) [41] form IR.

As discussed in § III-D, our approach utilizes a k-CFA call graph construction algorithm. To make our experiments tractable and to treat client-side API invocations as stream creations (since the focus of this work is on manipulation of client

TABLE III
EXPERIMENTAL RESULTS.

subject	KLOC	eps	k	str	rft	P1	P2	P3	t (m)
htm.java	41.14	21	4	34	10	0	10	0	1.85
JacpFX	23.79	195	4	4	3	3	0	0	2.31
jdp*	19.96	25	4	28	15	1	13	1	31.88
jdk8-exp*	3.43	134	4	26	4	0	4	0	0.78
jetty	354.48	106	4	21	7	3	4	0	17.85
jOOQ	154.01	43	4	5	1	0	1	0	12.94
koral	7.13	51	3	6	6	0	6	0	1.06
monads	1.01	47	2	1	1	0	1	0	0.05
retroλ	5.14	1	4	8	6	3	3	0	0.66
streamql	4.01	92	2	22	2	0	2	0	0.72
threeten	27.53	36	2	2	2	0	2	0	0.51
Total	641.65	751	4	157	57	10	46	1	70.60

* jdp is java-design-patterns and jdk8-exp is jdk8-experiments.

code), we made k an input parameter to our analysis (with $k = 2$ being the default as it is the minimum k value to consider client-code) for methods returning streams and $k = 1$ elsewhere. Recall that k amounts to the call string length in which to approximate object instances, thus, $k = 1$ would consider constructor calls as object creation locations, while $k = 2$ would consider calls to methods calling constructors as ("client") object creation sites. The tool currently uses a heuristic to inform developers when k is too small via a precondition failure. It does so by checking that call strings include at least one client method starting from the constructor call site. Future work involves automatically determining an optimal k, perhaps via stochastic optimization. The call graph used in the typestate analysis is pruned by removing nodes that do not have reaching stream definitions.

B. Experimental Evaluation

Our evaluation involved studying 11 open source Java applications and libraries of varying size and domain (table III). Subjects were also chosen such that they are using Java $>= 8$ and have at least one stream declaration (i.e., a call to a stream API) that is reachable from an entry point (i.e., a candidate stream). Column **KLOC** denotes the thousands of source lines of code, which ranges from ~1K for monads to ~354K for jetty. Column **eps** is the number of entry points. For non-library subjects, all main methods were chosen, otherwise, all unit test methods were chosen as entry points. Column **k** is the maximum k value used (see § IV-B1). Subjects compiled correctly and had identical unit test (27,955; mostly from jetty) results and compiler warnings before and after the refactoring.

The analysis was executed on an Intel Xeon E5 machine with 16 cores and 30GB RAM and a 25GB maximum heap size. Column **tm (m)** is the running time in minutes, averaging ~6.602 secs/KLOC. An examination of three of the subjects revealed that over 80% of the run time was for the typestate analysis, which is performed by SAFE. This analysis incorporates aliasing information and can be lengthy for larger applications. However, since our approach is automated, it can be executed on a nightly basis or before major releases.

1) Setting k for the k-CFA: As discussed in § III-D, our approach takes as input a maximum call string length parameter k, which is used to construct the call graph using nCFA. Each call graph node is associated with a *context*, which, in our case, is the call string. This allows our analysis to approximate stream object creation in the *client* code rather than in the framework, where the stream objects are instantiated. Otherwise, multiple calls to the same API methods that create streams would be considered as creating one new stream.

TABLE IV
REFACTORING FAILURES.

failure	pc	cnt
F1. InconsistentPossibleExecutionModes		1
F2. NoStatefulIntermediateOperations	P5	1
F3. NonDeterminableReductionOrdering		5
F4. NoTerminalOperations		13
F5. CurrentlyNotHandled		16
F6. ReduceOrderingMatters	P3	19
F7. HasSideEffects	P1	4
	P2	41
Total		100

During our experiments, a default k value of 2 was used. This is the minimum k value that can be used to distinguish client code from framework stream creation. However, depending on which stream framework methods are utilized in a particular project, this value may be insufficient. We detect this situation via a heuristic of examining the call string and determining whether any client code exists. If not, k may be too small.

Setting k constitutes a trade-off. A k that is too small will produce correct results but may miss streams. A larger k may enable the tool to detect and subsequently analyze more streams but may increase run time. Thus, an optimal k value can be project-specific. In our experiments, however, we determined k empirically based on a balance between run time and the ratio between total (syntactically available) streams and candidate streams (i.e., those detected by the typestate analysis). Notwithstanding, in keeping k between 2 and 4 (cf. table III), good results and reasonable runtime were observed. Thus, it was not difficult to find an "effective" k.

2) Intelligent Parallelization: Streams are still relatively new, and, as they grow in popularity, we expect to see them used more widely. Nevertheless, we analyzed 157 (origin) streams reachable from entry points (column **str**) across 11 subjects. Of those, we automatically refactored ~36.31% (column **rft** for *refactorable*) despite being highly conservative. These streams are the ones that have passed all preconditions; those not passing preconditions were not transformed (cf. table IV).

Columns **P1–3** are the streams passing the corresponding preconditions (cf. tables I and II). Columns **P4–5** have been omitted as all of their values are 0. The number of transformations can be derived from these columns as preconditions are associated with transformations, amounting to $10+46+(1*2) = 58$.

3) Refactoring Failures: Table IV categorizes reasons why streams could not be refactored (column **failure**), some of which correspond directly to preconditions (column **pc**). Column **cnt** depicts the count of failures in the respective category and further categorized by precondition, if applicable. Significant reasons streams were not refactorable include λ-expression side-effects (F7, 45%) and that the reduction ordering is preserved by the target collection (19%, c.f. § II).

Some of the refactoring failures were due to cases currently not handled by our tool (F5), which are rooted in implemen-

TABLE V
AVERAGE RUN TIMES OF JMH BENCHMARKS.

#	benchmark	orig (s/op)	refact (s/op)	su
1	shouldRetrieveChildren	0.011 (0.001)	0.002 (0.000)	6.57
2	shouldConstructCar	0.011 (0.001)	0.001 (0.000)	8.22
3	addingShouldResultInFailure	0.014 (0.000)	0.004 (0.000)	3.78
4	deletionShouldBeSuccess	0.013 (0.000)	0.003 (0.000)	3.82
5	addingShouldResultInSuccess	0.027 (0.000)	0.005 (0.000)	5.08
6	deletionShouldBeFailure	0.014 (0.000)	0.004 (0.000)	3.90
7	specification.AppTest.test	12.666 (5.961)	12.258 (1.880)	1.03
8	CoffeeMakingTaskTest.testId	0.681 (0.065)	0.469 (0.009)	1.45
9	PotatoPeelingTaskTest.testId	0.676 (0.062)	0.465 (0.008)	1.45
10	SpatialPoolerLocalInhibition	1.580 (0.168)	1.396 (0.029)	1.13
11	TemporalMemory	0.013 (0.001)	0.006 (0.000)	1.97

tation details related to model differences between representations [28]. For example, streams declared inside inner (embedded) classes are problematic as such classes are part of the outer AST but the instruction-based IR is located elsewhere. Though we plan to develop more sophisticated mappings in the future, such failures only accounted for 16%. Other refactoring failures include F4, where stream processing does not end with a terminal operation in all possible executions. This amounts to "dead" code as any queued intermediate operations will never execute. F3 corresponds to the situation described in § III-E3b, F1 to the situation where execution modes are ambiguous on varying execution paths, and F2 means that the stream is already optimized.

4) Performance Evaluation: Many factors can influence performance, including dataset size, number of available cores, JVM and/or hardware optimizations, and other environmental activities. Nevertheless, we assess the performance impact of our refactoring. Although this assessment is focused on our specific refactoring and subject projects, in the general case, it has been shown that a similar refactoring done manually has improved performance by 50% on large datasets [42, Ch. 6].

a) Existing Benchmarks: We assessed the performance impact of our refactoring on the subjects listed in table III. One of the subjects, htm.java [43], has formal performance tests utilizing a standard performance test harness, namely, the Java Microbenchmark Harness (JMH) [44]. Using such a test harness is important in isolating causes for performance changes to the code changes themselves [42, Ch. 6.1]. As such, subjects with JMH tests will produce the best indicators of performance improvements. Two such tests were included in this subject.

b) Converted Benchmarks: Although the remainder of the subjects did not include formal performance tests, they did include a rich set of unit tests. For one subject, namely, java-design-patterns [45], we methodically transformed existing JUnit tests that covered the refactored code to proper JMH performance tests. This was accomplished by annotating existing `@Test` methods with `@Benchmark`, i.e., the annotation that specifies that a method is a JMH performance test. We also moved setup code to `@Before` methods, i.e., those that execute before each test, and annotated those with `@Setup`. This ensures that the test setup is not included in the performance assessment. Furthermore, we chose unit tests that did not overly involve I/O (e.g., database access) to minimize variability. In all, nine unit tests were converted to performance tests and made our changes available to the subject developers.

c) Augmenting Dataset Size: As all tests we designed for continuous integration (CI), they executed on a minimal amount of data. To exploit parallelism, however, we augmented test dataset sizes. For existing benchmarks, this was done under the guidance of the developers [46]. For the converted tests, we chose an N (dataset size) value that is consistent with that of the largest value used by Naftalin [42, Ch. 6]. In this instance, we preserved the original unit test assertions, which all passed. This ensures that, although N has increased, the spirit of the test, which may reflect a real-life scenario, remains intact.

d) Results: Table V reports the average run times of five runs in seconds per operation. Rows 1–9 are for java-design-patterns, while rows 10–11 are for htm.java; benchmark names have been shortened for brevity. Column **orig** is the original program, **refact** is the refactored program, and **su** is the speedup ($runtime_{old}/runtime_{new}$). Values associated with parentheses are averages, while the value in parenthesis is the corresponding standard deviation. The average speedup resulting from our refactoring is 3.49.

5) Discussion: The findings of Naftalin [42, Ch. 6] using a similar manual refactoring, that our tool was able to refactor 36.31% of candidate streams (table III), and the results of the JMH tests on the refactored code (table V) combine to form a reasonable motivation for using our approach in real-world situations. Moreover, this study gives us insight into how streams, and in a broader sense, concurrency, are used, which can be helpful to language designers, tool developers, and researchers.

As mentioned in § IV-B2, columns **P4–5** in table III all have 0 values. Interestingly, this means that no (already) parallel streams were refactored by our tool. Only two candidate streams, stemming from only a single subject, htm.java, were originally parallel. This may indicate that developers are either timid to use parallel streams because of side-effects, for example, or are (manually) unaware of when using parallel streams would improve performance [42]. This further motivates our approach for automated refactoring in this area.

From table IV, F6 and F7 accounted for the largest percentage of failures (64%). For the latter, this may indicate that despite that "many computations where one might be tempted to use side-effects can be more safely and efficiently expressed without side-effects" [5], in practice, this is either not the case or more developer education is necessary to avoid side-effects when using streams. This motivates future work in refactoring stream code to avoid side-effects if possible.

Imprecision is also a possibility as we are bound by the conservativeness of the underlying ModRef analysis provided by WALA. To investigate, we manually examined 45 side-effect failures and found 11 false positives. Several subject developers, on the other hand, confirmed correct refactorings, as discussed in § IV-B6. As for the former, a manual inspection of these sites may be necessary to confirm that ordering indeed must be preserved. If not, developers can rewrite the code (e.g., changing `forEachOrdered()` to `forEach()`) to exploit more parallelism opportunities.

The average speedup of 1.55 obtained from htm.java (benchmarks 10–11) most likely reflects the parallelism opportunities available in computationally intensive programs [47]. Benchmarks 1–6, which had good speedups as well, also mainly deal with data. Benchmark 7 had the smallest speedup at 1.03. The problem is that the refactored code appears in areas that "will not benefit from parallelism" [48], demonstrating a limitation of our approach that is rooted in its problem scope. Specifically, our tool locates sites where stream client code is safe to refactor and is possibly optimizable based on language semantics but does not assess optimizability based on input size/overhead trade-offs.

6) Pull Request Study: To assess our approach's usability, we also submitted several pull requests (patches) containing the results of our tool to the subject projects. As of this writing, eight requests were made, with three pending (e.g., [46]) and five rejected. One rejected request [48] is discussed in § IV-B5. Others (e.g., [45]) confirmed a correct refactoring but only wanted parallel streams when performance is an observed problem.

C. Threats to Validity

The subjects may not represent the stream client code usage. To mitigate this, subjects were chosen from diverse domains as well as sizes, as well as those used in previous studies (e.g., [49], [50]). Although java-design-patterns is artificial, it is a reference implementation similar to that of JHotDraw, which has been studied extensively (e.g., [51]).

Entry points may not be correct, which would affect which streams are deemed as candidates, as well as the performance assessment as there is a trade-off between scalability and number of entry points. Since standard entry points were chosen (see § IV-B), representing a super set of practically true entry points. For the performance test (see table V), the loads may not be representative of real-world usage. However, we conferred with developers regarding this when possible [46]. For the performance tests we manually generated from unit tests, a systematic approach to the generation was taken using the same parameters (N) on both the original and refactored versions.

V. Related Work

Automatic parallelization can occur on several levels, including the compiler [52], [53], run time [54], and source [17]. The general problem of full automatic parallelization by compilers is extremely complex and remains a grand challenge [55]. Many attempt to solve it in only certain contexts, e.g., for divide and conquer [56], recursive functions [57], distributed architectures [58], graphics processing [59], matrix manipulation [60], asking the developer for assistance [61], and speculative strategies [62]. Our approach focuses on MapReduce-style code over native data containers in a shared memory space using a mainstream programming languages, which may be more amenable to parallelization due to more explicit data dependencies [16]. Moreover, our approach can help detect when it is *not* advantageous to run code in parallel, and when unordering streams can possibly improve performance.

Techniques other than ours enhance the performance of streams as well. Hayashi *et al.* [63] develop a supervised machine-learning approach for building performance heuristics for mapping Java applications onto CPU/GPU accelerators via analyzing parallel streams. Ishizaki *et al.* [64] translate λ-expressions in parallel streams into GPU code and automatically generates run time calls that handle low-level operations. While all these approaches aim to improve performance, their input is streams that are *already* parallel. As such, developers must still *manually* identify and transform sequential streams. Nonetheless, these approaches may be used in conjunction with ours.

Harrison [65] develops an interprocedural analysis and automatic parallelization of Scheme programs. While Scheme is a multi-paradigm language, and shared memory is modeled, their transformations are more invasive and imperative-focused, involving such transformations as eliminating recursion and loop fusion. Nicolay *et al.* [66] have a similar aim but are focused on analyzing side-effects, whereas we analyze ordering constraints.

Many approaches use streams for other tasks or enhance streams in some way. Cheon *et al.* [67] use streams for JML specifications. Biboudis *et al.* [1] develop "extensible" pipelines that allow stream APIs to be extended without changing library code. Other languages, e.g., Scala [2], JavaScript [3], C# [4], also offer streaming APIs. While we focus on Java 8 streams, the concepts set forth here may be applicable to other situations, especially those involving statically-typed languages, and is a topic for future work.

Other approaches refactor programs to either utilize or enhance modern construct usage. Gyori *et al.* [16] refactor Java code to use λ-expressions instead of imperative-style loops. Tsantalis *et al.* [68] transform clones to λ-expressions. Khatchadourian and Masuhara [69] refactor skeletal implementations to default methods. Tip *et al.* [70] use type constraints to refactor class hierarchies, and Gravley and Lakhotia [71] and Khatchadourian [72] refactor programs to use enumerated types.

Typestate has been used to solve many problems. Mishne *et al.* [73] use typestate for code search over partial programs. Garcia *et al.* [74] integrate typestate as a first-class citizen in a programming language. Padovani [75] extends typestate oriented programming (TSOP) for concurrent programming. Other approaches have also used hybrid typestate analyses. Bodden [76], for instance, combines typestate with residual monitors to signal property violations at run time, while Garcia *et al.* [74] also make use of run time checks via gradual typing [75].

VI. Conclusion & Future Work

Our automated refactoring approach "intelligently" optimizes Java 8 stream code It automatically deems when it is safe and possibly advantageous to run stream code either sequentially or in parallel and unorder streams. The approach was implemented as an Eclipse plug-in and evaluated on 11 open source programs, where 57 of 157 candidate streams (36.31%) were refactored. A performance analysis indicated an average speedup of 3.49.

In the future, we plan to handle several issues between Eclipse and WALA models and incorporate more kinds of (complex) reductions like those involving maps, as well as look into approximations to combat the problems set forth by Chen *et al.* [20]. Approximating SIOs may also involve heuristics, e.g., analysis of API documentation. Lastly, we will explore applicability to other streaming frameworks and languages.

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